

Computing Inconsistency Measures Under Differential Privacy

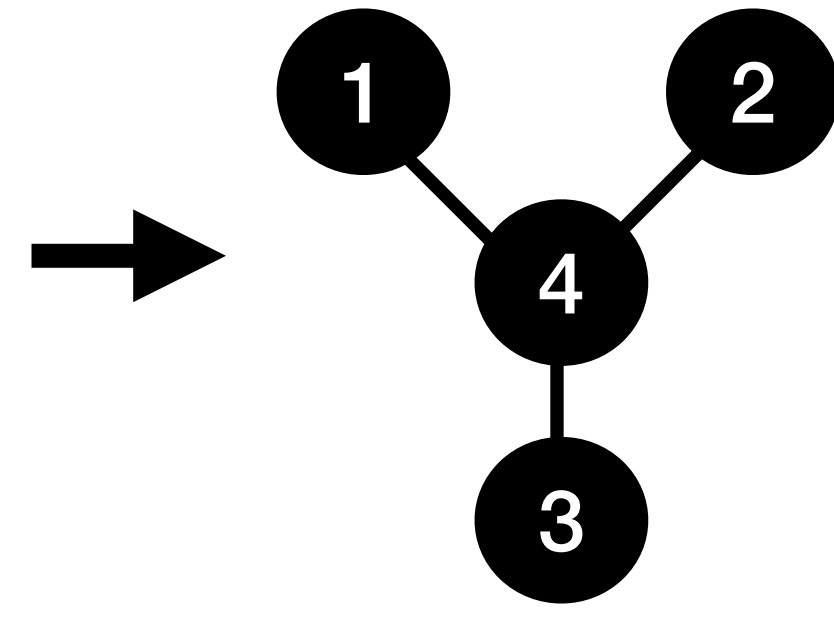
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Background

Denial Constraints and Conflict Graphs

ID	Capital	Country
1	Ottawa	Canada
2	Ottawa	Canada
3	Ottawa	Canada
4	Ottawa	Kanada



σ : Capital $\rightarrow$ Country

“country of two tuples must be the same if their capital is the same”

Inconsistency Measures [1]

Given, dataset D and constraint set Σ , inconsistency measures are of the form $I(D, \Sigma) \rightarrow \mathbb{R}$:

1. Drastic Measure $I_D(G)$ = existence of an edge ✗
2. Minimal inconsistency measure $I_M(G)$ = number of edges ✓
3. Problematic measure $I_P(G)$ = number of vertices with positive degree ✓
4. Maximal consistency measure $I_{MC}(G)$ = number of maximal independent sets ✗
5. Optimal repair measure $I_R(G)$ = minimum vertex cover size ✓

✓ = our work

Differential Privacy [2]

A randomized algorithm $A: \mathcal{G} \rightarrow \mathcal{R}$ satisfies ϵ -differential privacy (DP) if for any two adjacent graphs $G, G' \in \mathcal{G}$ that differ in a node and for any subset of outputs $o \subseteq \mathcal{R}$ it holds that :

$$\Pr[A(G) \in o] \leq e^\epsilon \Pr[A(G') \in o]$$

Experiments

Setup

- Inconsistency: Random typo to 1% rows
- Privacy: $\epsilon = 1$
- Measure: True vs private by adding one typo at a time

Datasets

- Five real-world datasets with varying conflict graph densities
- We experiment on subset of 10k and repeat for 10 times and average

Dataset	#Tuples	#Attributes	#Constraints	Graph density
Adult	32k	15	3	9635
Flight	500k	20	13	1520
Hospital	114k	15	7	793
Stock	122k	7	1	1
Tax	1M	15	9	373

References:

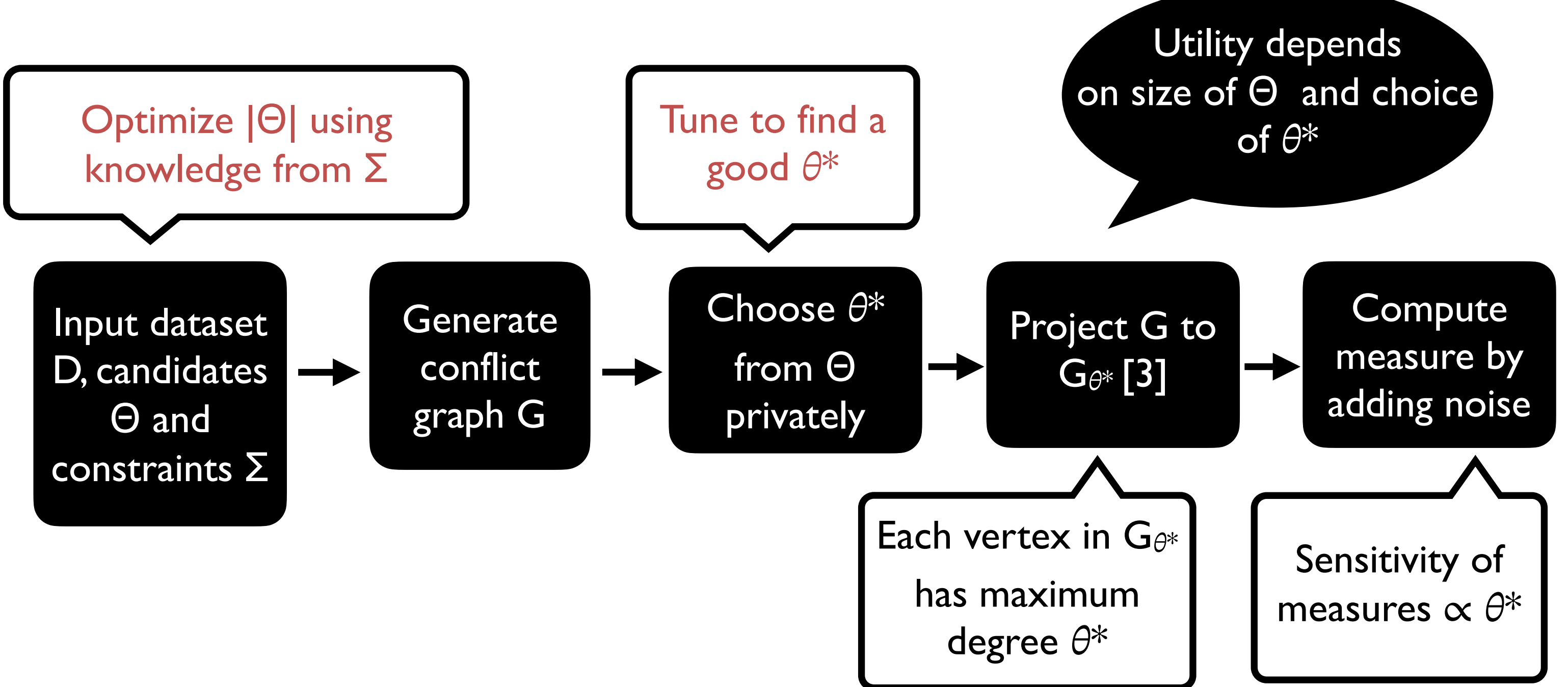
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3. Day, Wei-Yen, et al. "Publishing graph degree distribution with node differential privacy." *International Conference on Management of Data* 2016.
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Differentially Private Inconsistency Measures

Challenges

1. Computational hardness: Minimum vertex cover (I_R) and #maximal independent sets (I_{MC}) are NP hard problems
2. High sensitivity: Maximum change in output when G is replaced by G' is $\mathcal{O}(n)$, where n is the number of nodes

Graph projection for I_P and I_M [$\theta^* \ll \mathcal{O}(n)$]



DP vertex cover size for I_R [$2 \ll \mathcal{O}(n)$]

We analyse the 2-approximate vertex cover size algorithm and show that it has sensitivity of 2.

